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# Policy and Value Transfer in Lifelong Reinforcement Learning (Appendix)

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David Abel <sup>†1</sup> Yuu Jinnai <sup>†1</sup> Yue Guo <sup>1</sup> George Konidaris <sup>1</sup> Michael L. Littman <sup>1</sup>

We here include proofs and a visual of the Octogrid domain.

## A. Proofs

**Theorem 3.2.** For a distribution of MDPs with  $R \sim D$ ,

$$\mathbb{E}_{M \in \mathcal{M}} [V_M^{\pi^{*avg}}(s)] \geq \max_{M \in \mathcal{M}} \Pr(M) V_M^*(s).$$

*Proof.* Ramachandran Amir (2007) also showed that the value function  $V_{avg}^{\pi}$  of an average MDP is the weighted average of the MDPs in the distribution,

$$V_{avg}^{\pi}(s) = \sum_{M \in \mathcal{M}} \Pr(M) V_M^{\pi}(s). \quad (1)$$

Thus,

$$\begin{aligned} \mathbb{E}_{M \in \mathcal{M}} [V_M^{\pi^{*avg}}(s)] &= \sum_{M \in \mathcal{M}} \Pr(M) V_M^{\pi^{*avg}}(s) \\ &= V_{avg}^{\pi^{*avg}}(s) \\ &= \max_{\pi} V_{avg}^{\pi}(s) \\ &= \max_{\pi} \sum_{M \in \mathcal{M}} \Pr(M) V_M^{\pi}(s) \\ &\geq \max_{\pi} \max_{M \in \mathcal{M}} \Pr(M) V_M^{\pi}(s) \\ &= \max_{M \in \mathcal{M}} \Pr(M) \max_{\pi} V_M^{\pi}(s) \\ &= \max_{M \in \mathcal{M}} \Pr(M) V_M^*(s). \end{aligned}$$

Since we assume  $\mathcal{R}(s, a) \geq 0$  for all  $s, a$ , we infer that  $\sum_{M \in \mathcal{M}} \Pr(M) V_M^{\pi}(s) \geq \max_{M \in \mathcal{M}} \Pr(M) V_M^{\pi}(s)$ , thus concluding the proof.  $\square$

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**Corollary 3.2.1.** The bound in Theorem 3.2 is tight.

*Proof.* Next we the bound is by an example MDP distribution shown in Figure 1.

In the MDP  $i$  the agent gets a reward if it executes  $a_i$  in MDP  $i$ :

$$R_M(s_0, a_i) = \begin{cases} 1 & M = i \\ 0 & \text{otherwise} \end{cases}$$

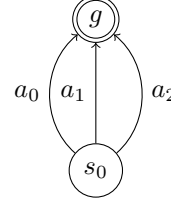


Figure 1: An example of a MDP which an average MDP solution returns a lower bound value.

In this distribution of MDPs, the optimal agent always gets reward of 1 where as the optimal average MDP agent gets  $\max_{M \in \mathcal{M}} \Pr(M)$  reward on average. In this setting,  $V^{\pi^{*avg}}(s) = \max_{M \in \mathcal{M}} \Pr(M) V_M^*(s)$ . Thus the bound is tight.  $\square$

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**Corollary 3.4.** For the  $G \sim D$  setting,

$$\begin{aligned} \mathbb{E}_{M \in \mathcal{M}} [V_M^{\pi^{*avg}}(s)] &\geq \min_{M \in \mathcal{M}} \Pr(M) \max_{M' \in \mathcal{M}} \Pr(M') V_{M'}^*(s). \end{aligned}$$

*Proof.* We first leverage the following lemma:

**Lemma 3.4.1.**

$$\begin{aligned} \max_{M \in \mathcal{M}} \Pr(M) V_M^{\pi}(s) &\leq V_{avg}^{\pi}(s) \\ &\leq \sum_{M \in \mathcal{M}} \Pr(M) V_M^{\pi}(s) / \min_{M' \in \mathcal{M}} \Pr(M') \end{aligned}$$

(*Proof sketch for lower bound*): Let an MDP  $M'$  be the same MDP as  $M$  except it transits to a terminal state from goal nodes (and acquires a reward) by probability of  $\Pr(M)$  instead of probability of 1. The value  $V_{M'}^{\pi}(s)$  of state  $s$  in  $M'$  is at least as large as  $\Pr(M) V_M^{\pi}(s)$ . Thus, the value of state  $s$  in  $M'$  is lower than or equal to that in the average MDP as it reaches the goal less frequently.  $V_{M'}^{\pi}(s)$  is smaller that or equal to  $V_{avg}^{\pi}(s)$  as the average MDP has larger or equal probability of reaching the terminal state. Thus, for any  $M \in \mathcal{M}$ :

$$V_{avg}^{\pi}(s) \geq V_{M'}^{\pi}(s) \geq \Pr(M) V_M^{\pi}(s).$$

(Proof sketch for upper bound):

$$\begin{aligned} V_{avg}^\pi(s) &\leq \sum_{M \in \mathcal{M}} V_M^\pi(s) \\ &\leq \sum_{M \in \mathcal{M}} \Pr(M) V_M^\pi(s) / \min_{M' \in \mathcal{M}} \Pr(M'). \end{aligned}$$

Now, we turn to the theorem.

$$\begin{aligned} \mathbb{E}_{M \in \mathcal{M}} [V_M^{\pi_{avg}^*}(s)] &= \sum_{M \in \mathcal{M}} \Pr(M) V_M^{\pi_{avg}^*}(s) \\ &\geq \min_{M \in \mathcal{M}} \Pr(M) V_{avg}^{\pi_{avg}^*}(s) \\ &= \min_{M \in \mathcal{M}} \Pr(M) \max_{\pi} V_{avg}^\pi(s) \\ &\geq \min_{M \in \mathcal{M}} \Pr(M) \max_{\pi} \max_{M' \in \mathcal{M}} \Pr(M') V_{M'}^\pi(s) \\ &= \min_{M \in \mathcal{M}} \Pr(M) \max_{M' \in \mathcal{M}} \Pr(M') V_{M'}^*(s). \quad \square \end{aligned}$$

**Theorem 3.8.** Suppose  $\mathcal{A}$  is an algorithm that produces  $\varepsilon$  accurate  $Q$  functions for a subset of the state action space given an MDP  $M$ , an initial state  $s_0$ , and a horizon  $H$ . For a given  $\delta \in (0, 1]$ , after

$$t \geq \frac{\ln(\delta)}{\ln(1 - p_{min})}, \quad (2)$$

sampled MDPs, for  $p_{min} = \min_{M \in \mathcal{M}} \Pr(M)$ , the updating-max shaping method will return a shaped  $Q$ -function  $\hat{Q}_{max}$  such that for all state action pairs  $(s, a)$ :

$$\hat{Q}_{max}(s, a) \geq \max_M Q_M^*(s, a), \quad (3)$$

with probability  $1 - \delta$ .

*Proof.* Consider an arbitrary state action pair  $(s, a)$ .

After  $t$  samples, we choose:

$$\hat{Q}_{max}^*(s, a) \triangleq \max_M \hat{Q}_M^*(s, a). \quad (4)$$

After  $t$  samples, we let the following event define a mistake:

$$\hat{Q}_{max}^*(s, a) < \max_M Q_M^*(s, a). \quad (5)$$

First, we suppose that for each of sampled MDP  $M$ , our learning algorithm computes a *partial* but nearly *accurate*  $Q$ -function. That is, for some small  $\varepsilon$ :

$$\hat{Q}_M^*(s, a) = \begin{cases} Q_M^*(s, a) \pm \varepsilon & c(s, a) \geq m \\ \text{VMAX} & \text{otherwise} \end{cases} \quad (6)$$

That is, letting  $c(s, a)$  denote the number of times  $a$  was executed in  $s$ : any state action pairs that were visited sufficiently often (more than  $m$  for some chosen  $m \ll H$ )

result in an  $\varepsilon$ -accurate  $Q$  function. Otherwise, the algorithm returns VMAX.

Under these conditions, for a given state action pair, surely, for any MDP seen during the  $t$  samples  $M_i$ :

$$\hat{Q}_{max}^*(s, a) \geq \max_{M \in \mathcal{M}_{seen}} Q_M^*(s, a) \quad (7)$$

Therefore, the mistake event defined by Equation 5 only occurs when we *miss* an MDP in the distribution that has a higher  $Q^*(s, a)$  than our estimate. We assume that the distribution has a lower bound on MDP probability:

$$p_{min} \triangleq \min_{M \in \mathcal{M}} \Pr(M). \quad (8)$$

Accordingly, we upper bound the mistake probability according to the probability that no such MDP was sampled over  $t$  samples, captured by the cumulative geometric distribution:

$$1 - (1 - p_{min})^m \geq 1 - \delta. \quad (9)$$

Simplifying:

$$\begin{aligned} 1 + \delta &\geq 1 + (1 - p_{min})^t \\ \ln(\delta) &\geq \ln(1 - p_{min}) \cdot t \\ \frac{\ln(\delta)}{\ln(1 - p_{min})} &\leq t \end{aligned}$$

Therefore, after

$$t \geq \frac{\ln(\delta)}{\ln(1 - p_{min})}, \quad (10)$$

sampled MDP we will have seen all MDPs in the distribution with high probability.  $\square$

## B. Octogrid

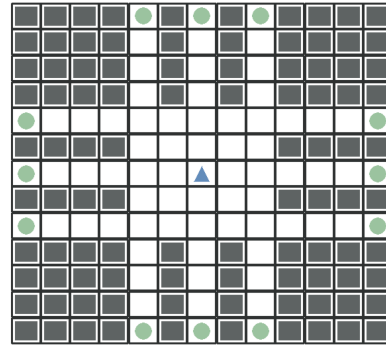


Figure 2: The Octogrid task distribution. The goal appears in exactly one of the 12 green circles chosen uniformly at random, with the agent starting in the center at the triangle.